Conditions for regularity and for 2-connectivity of Toeplitz graphs

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Abstract

Let $1 \leq t_1 < t_2 < \cdots < t_k \leq n$. A Toeplitz graph G = (V, E) denoted by $T_n(t_1, \ldots, t_k)$ is a graph where $V = \{1, \ldots, n\}$ and $E = \{(i, j) \mid |i - j| \in \{t_1, \ldots, t_k\}\}$. In this paper, we classify all regular Toeplitz graphs. Here, we present some conditions under which a Toeplitz graph has no cut-edge and cut-vertex.

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1. Introduction

A Toeplitz matrix is named after Otto Toeplitz (1881-1940) which is an $n \times n$ matrix $A = (a_{ij})$ such that for each i and j, $1 \leq i, j \leq n - 1$, $a_{ij} = a_{(i+1)(j+1)}$. Toeplitz (0, 1)-matrices are precisely those matrices that all diagonals parallel to main diagonal has constant values. Thus, Toeplitz matrices are determined by its first row and column. Let n, t_1, \ldots, t_k be distinct positive integers such that $1 \leq t_1 < t_2 < \cdots < t_k < n$. A Toeplitz graph is denoted by $T_n\langle t_1, \ldots, t_k \rangle = (V, E)$ where n is the number of vertices, $V = \{1, \ldots, n\}$ and $E = \{\{i, j\} \mid |i - j| \in \{t_1, \ldots, t_k\}\}$. The name of this class of graphs is due to the fact that their adjacency matrices is a



Figure 1: The Toeplitz graph $T_7 < 3, 4, 5 >$

Toeplitz (0, 1)-matrix. For example, see the graph $T_7 < 3, 4, 5 >$, shown in Figure 1. Moreover, the number of edges in the Toeplitz graph $T_n(t_1,\ldots,t_k)$ is equal to $\sum_{i=1}^{k} (n - t_i)$, see [2]. Properties of Toeplitz graphs, such as biparticity, planarity, colourabil-

ity and Hamiltonicity have been studied in [1]-[12].

Let G be a graph with the vertex set V(G) and the edge set E(G). We denote the degree of v in G by d(v) and minimum degree of G is denoted by $\delta(G)$. A graph G is r-regular if d(v) = r, for all $v \in V(G)$. The set of neighbors of v in graph G is denoted by $N_G(v)$ or simply N(v). For a set $S \subseteq V$, its open neighborhood is $N(S) = \bigcup_{v \in S} N(v)$. If G is a graph and $S \subseteq V(G)$, the induced subgraph on S is denoted by G[S]. The cycle of order n is denoted by C_n . A graph G is said to be k-vertex-connected (or k-connected) if the graph remains connected after deleting any fewer than k vertices from the graph. A *cut vertex* of a graph G is a vertex v whose deletion along with incident edges results in a graph with more components than the original graph. A graph is k-edge-connected if it remains connected whenever any fewer than k edges are removed. A cut edge of a graph Gis an edge e whose deletion results in a graph with more components than the original graph.

2. Vertex Degrees in Toeplitz Graphs

In this section we present a result on the degree of vertices of Toeplitz graphs, and some results on the minimum degree of Toeplitz graph and characterize *r*-regular Toeplitz graphs.

Lemma 1 Let $G = T_n < t >$. For even integer n and each $i \in V(G)$, d(i) = d(n-i+1). For odd integer n, d(i) = d(n-i+1) for every $i \in V(G) \setminus \{ \lceil \frac{n}{2} \rceil \}.$

Proof. Case 1. Let $t = \lfloor \frac{n}{2} \rfloor$.

For even n, we have $E(G) = \{(i, t+i); 1 \le i \le t\}$, which clearly shows that d(i) = 1 for each $i \in V(G)$. Hence, for every vertex i in G, d(i) = d(n-i+1). For odd n, we have $E(G) = \{(i, t+i); 1 \le i \le t\} \cup (t+1, n)$. Here, since t+1 is the only vertex with degree two, d(i) = 1, for every $i \in V(G) \setminus \{ \lfloor \frac{n}{2} \rfloor \}$. Hence d(i) = d(n - i + 1), for every $i \in V(G) \setminus \lfloor \frac{n}{2} \rfloor$.

Case 2. Let $t > \lfloor \frac{n}{2} \rfloor$.

In this case we have $E(G) = \{(i, t+i); 1 \le i \le n-t\}$, which clearly shows that

$$d(i) = \begin{cases} 1, & for \ i \in \{1, 2, \dots, n-t\} \cup \{t+1, t+2, \dots, n\} \\ \\ 0, & for \ i \in \{n-t+1, n-t+2, \dots, t\}. \end{cases}$$

and it can be easily seen that for even n, d(i) = d(n - i + 1), for each $i \in V(G)$. Also for odd n, d(i) = d(n - i + 1), for all $i \in V(G) \setminus \{\lceil \frac{n}{2} \rceil\}$, because $d(\lceil \frac{n}{2} \rceil) = d(n - \lceil \frac{n}{2} \rceil + 1) = d(\lfloor \frac{n}{2} \rfloor + 1) = d(\lceil \frac{n}{2} \rceil)$ and there are odd number of vertices with degree zero.

Case 3. Let $t < \lfloor \frac{n}{2} \rfloor$.

In this case G consists of the paths $\{(i, i+t, i+2t, \dots, m); 1 \le i \le t\}$, where m is the greatest integer less than or equal to n such that $m \cong i \pmod{t}$. So $d(i) \in \{1, 2\}$, for $1 \le i \le n$. Now, by induction on $i, 1 \le i \le \lfloor \frac{n}{2} \rfloor$, we will show that d(i) = d(n - i + 1). For i = 1 the result is true, because d(1) = 1 = d(n). Suppose that the result holds for i = k, i.e., d(k) = d(n-k+1). We will now show that d(k+1) = d(n-k). To the contrary, suppose $d(k+1) \neq d(n-k)$. Then, without loss of generality, d(k+1) = 2 and d(n-k) = 1. Since d(k+1) = 1 and $t < \lfloor \frac{n}{2} \rfloor$, (k+1)+t < nand (k+1) - t < 1. Together, these imply $k + t < n - \overline{1}$. And k - t < 0and therefore d(k) = 1. By induction hypothesis, d(n - k + 1) = 1. It now follows that (n-k+1)-t = n+1-(k+t) > 1 and (using the fact that k+t < n (n-k+1)+t > n. Together, these imply n-(k+t) > 0 and (n-k)+t > n-1. Since d(n-k) = 2, it must be the case that n > k+tand (n-k) + t = n. Thus n > 2t, contradicting that $t < \lfloor n/2 \rfloor$, proving the Lemma." \square

Proposition 1 Let $G = T_n < t_1, \ldots, t_k >$. If *n* is even, then d(i) = d(n - i + 1) for each $i \in V(G)$. Also, if *n* is odd, then for each $i \in V(G) \setminus \{\lceil \frac{n}{2} \rceil\}$, d(i) = d(n - i + 1).

Proof. We prove the proposition by induction on k. Lemma 1 shows that the result is true for k = 1. Suppose that the result is true for k = s, i.e., for $T_n < t_1, t_2, \ldots, t_s >$. Now we show that the result is true for k = s + 1. We can easily see that

$$T_n < t_1, \dots, t_{s+1} >= T_n < t_1 > \cup T_n < t_2, t_3, \dots, t_{s+1} >$$

By induction hypothesis the result is true for both $T_n < t_1 >$ and $T_n < t_2, t_3, \ldots, t_{s+1} >$. Clearly, $E(T_n < t_1 >) \cap E(T_n < t_2, t_3, \ldots, t_{s+1} >) = \emptyset$, which implies that the result is true for $T_n < t_1 > \cup T_n < t_2, t_3, \ldots, t_{s+1} >$, which completes the proof.

Lemma 2 For Toeplitz graph $G = T_n < t_1, \ldots, t_k >$, $\delta(G) = 0$ if and only if $t_1 \geq \left\lceil \frac{n+1}{2} \right\rceil$.

Proof. First assume that $\delta(G) = 0$. Let $i \in V(G)$ and d(i) = 0. Clearly, $i + t_1 \ge n + 1$. Thus $2t_1 \ge n + 1$ and so $t_1 \ge \lceil \frac{n+1}{2} \rceil$. Conversely, suppose that $t_1 \ge \lceil \frac{n+1}{2} \rceil$. If *n* is even, then $d(\frac{n}{2}) = 0$ and if *n*

is odd then $d(\frac{n+1}{2}) = 0$. Hence $\delta(\tilde{G}) = 0$. \square

Lemma 3 Let $G = T_n < t_1, ..., t_k > and n \ge 2$. Then $\delta(G) = 1$ if and only if $n < t_1 + t_2$ and $t_1 < \lceil \frac{n+1}{2} \rceil$.

Proof. First assume that $\delta(G) = 1$. Therefore, there exists $i \in V(G)$ such that d(i) = 1. So either $\{i, i - t_1\}$ or $\{i, i + t_1\} \in E(G)$. Without loss of generality, $\{i, i - t_1\} \in E(G)$. Since $\{i, i - t_2\}, \{i, i + t_1\} \notin E(G), i \le t_2$ and $n+1 \le i+t_1$. Thus $n+1 \le t_1+t_2$. By Lemma 2, $t_1 < \lceil \frac{n+1}{2} \rceil$.

Conversely, suppose that $n < t_1 + t_2$ and $t_1 < \lceil \frac{n+1}{2} \rceil$. The later one implies that $\delta(G) \geq 1$. Now, we show that G has at least one vertex of degree 1. If n is even, then by assumption, we have $2t_1 \leq n$ and so $\{t_1, 2t_1\} \in E(G)$. On the other hand, clearly for integer s, $1 \le s \le k$, $\{t_1, t_1 - t_s\} \notin E(G)$. Since $n < t_1 + t_2$, for integer $r, 2 \le r \le k$, $\{t_1, t_1 + t_r\} \notin E(G)$. Thus $d(t_1) = 1$. Similarly, if n is odd and $t_1 < \frac{n+1}{2}$, then $d(t_1) = 1$.

Now, we have the following corollary.

Corollary 1 Suppose that $G = T_n \langle t_1, \ldots, t_k \rangle$ and $n \geq 2$. Then $\delta(G) \geq 2$ if and only if $n \ge t_1 + t_2$ and $t_1 < \lceil \frac{n+1}{2} \rceil$.

Theorem 1.[7] A Toeplitz graph $T_n < t_1, \ldots, t_k > with t_1 + t_k \leq n + 1$ and $gcd(t_1, \ldots, t_k) = 1$ is a connected graph.

For k = 2, Theorem 1 states $G = T_{t_1+t_2} \langle t_1, t_2 \rangle$ is connected if $gcd(t_1, t_2) =$ 1, in Proposition 2, we show the graph is d-connected if $gcd(t_1, t_2) = d$.

Proposition 2 If $G = T_{t_1+t_2}\langle t_1, t_2 \rangle$ and $gcd(t_1, t_2) = d$, then G is the union of disjoint cycles, $G_0, G_1, \ldots, G_{d-1}$, and

$$V(G_i) = \{dj - i \mid 1 \le j \le \frac{t_1 + t_2}{d}\},\$$
$$E(G_i) = \{\{dj_1 - i, dj_2 - i\} \mid \{j_1, j_2\} \in E(T_{\frac{t_1 + t_2}{d}} \langle \frac{t_1}{d}, \frac{t_2}{d} \rangle)\},\$$
for $i = 0, 1, \dots, d - 1.$

Proof. Let $H = G_0 \cup G_1 \cup \cdots \cup G_{d-1}$. We claim that H is a subgraph of G. Since $d \ge 1$, for each $0 \le i \le d-1$ and $1 \le j \le \frac{t_1+t_2}{d}$, we have $1 \le dj-i \le t_1+t_2$. Thus $V(H) \subseteq V(G)$. Now, we show that $E(G_i) \subseteq E(G)$. Consider the edge $\{dj_1 - i, dj_2 - i\} \in E(G_i)$. Since $|j_1 - j_2| \in \{\frac{t_1}{d}, \frac{t_2}{d}\}$ and $|dj_1 - i - (dj_2 - i)| \in \{t_1, t_2\}$, therefore $\{dj_1 - i, dj_2 - i\} \in E(G)$. Hence $E(H) \subseteq E(G)$ and we have $H \subseteq G$. Now, we prove that $E(G) \subseteq E(H)$. Suppose that $\{r, s\} \in E(G)$. Let $r \equiv -i_1 \pmod{d}$ and $s \equiv -i_2 \pmod{d}$ for $0 \le i_1, i_2 \le d-1$. Since $\{r, s\} \in E(G), |r-s| \in \{t_1, t_2\}$. So $|r-s| \equiv 0 \pmod{d}$. Therefore, $|i_1 - i_2| \equiv 0 \pmod{d}$. Since $0 \le i_1, i_2 \le d-1$, $i_1 = i_2$. Assume that $i_1 = i_2 = i$. Thus, $r = dj_1 - i$ and $s = dj_2 - i$. So $j_1 = \frac{r+i}{d}$ and $j_2 = \frac{s+i}{d}$. Since $r, s \le t_1 + t_2, 0 \le \frac{i}{d} < 1$ and $j_1, j_2 \le \frac{t_1+t_2}{d}$. Hence, $\{r, s\} \in E(H)$, and we have $E(G) \subseteq E(H)$. Let $x \in V(G_a) \cap V(G_b)$ and $0 \le a < b \le d-1$. Therefore for some j_1 and $j_2, 1 \le j_1, j_2 \le \frac{t_1+t_2}{d}$, $x = dj_1 - a$ and $x = dj_2 - b$. So $a \equiv b \pmod{d}$. Since $0 \le a < b \le d-1$. Hence,

$$|V(H)| = \sum_{i=0}^{d-1} |V(G_i)| = d \frac{t_1 + t_2}{d} = t_1 + t_2 = |V(G)|.$$

Hence G = H.

Now, we have the following corollary.

Corollary 2 If $G = T_n \langle t_1, t_2 \rangle$ and $n \leq t_1 + t_2$, then G is a disjoint union of paths and cycles.

Now, we consider the Toeplitz graphs with $n \ge t_1 + t_2$ vertices.

Proposition 3 If $G = T_n \langle t_1, t_2 \rangle$, $n \ge t_1 + t_2$ and $gcd(t_1, t_2) = d$, then G is the disjoint union of $G[V_i]$, $0 \le i \le d-1$, where

$$V_i = \{k(t_1 + t_2) + dj - i \le n \mid 0 \le k \le \lfloor \frac{n}{t_1 + t_2} \rfloor \text{ and } 1 \le j \le \frac{t_1 + t_2}{d}\}$$

for each $i, 0 \leq i \leq d-1$.

Proof. Let $V_i = \{k(t_1 + t_2) + dj - i \le n \mid 0 \le k \le \lfloor \frac{n}{t_1 + t_2} \rfloor$ and $1 \le j \le \frac{t_1 + t_2}{d}\}$; where $d = gcd(t_1, t_2)$ and $0 \le i \le d - 1$.

Assume that $V_a \cap V_b \neq \emptyset$, for $0 \le a < b \le d - 1$. Then let $x \in V_a \cap V_b$. So $x = k_1(t_1 + t_2) + dj_1 - a = k_2(t_1 + t_2) + dj_2 - b$ for some k_1, k_2, j_1 and j_2 such that $0 \le k_1, k_2 \le \lfloor \frac{n}{t_1 + t_2} \rfloor$, $1 \le j_1, j_2 \le \frac{t_1 + t_2}{d}$. So $a \equiv b \pmod{d}$. Since $0 \le a, b \le d - 1, a = b$. Hence a contradiction.

Now, suppose that $\{r, s\} \in E(G)$ and s < r. Let $r = k_1(t_1 + t_2) + dj_1 - dj_1$

 i_1 and $s = k_2(t_1 + t_2) + dj_2 - i_2$, where $0 \le i_1, i_2 \le d - 1$. So $(k_1 - k_2)(t_1 + t_2) + d(j_1 - j_2) - (i_1 - i_2) \in \{t_1, t_2\}$. Since $gcd(t_1, t_2) = d$, $i_1 \equiv i_2$ (mod d). Since $0 \le i_1, i_2 \le d - 1$, $i_1 = i_2$ and consequently $\{r, s\} \in E(G[V_{i_1}])$. Thus $G[V_0], \ldots, G[V_{d-1}]$ are mutually disjoint subgraphs and $E(G) \subseteq \bigcup_{i=0}^{d-1} E(G[V_i])$.

It is straightforward to check that $V(G) = \bigcup_{i=0}^{d-1} V_i$. Therefore, the proof is complete.

Corollary 3 Let $G = T_{n-1}\langle t_1, \ldots, t_k \rangle$. If $n \ge t_1 + t_k$, then $\{n - t_1, n - t_2, \ldots, n - t_k\} \subseteq V(G)$ is contained in one of the components of G.

Proof. First, we prove that $n - t_1$ and $n - t_2$ are in the same component of $T_{n-1}\langle t_1, t_2 \rangle$. Consider the notations given in the Proposition 3. Suppose that $n - t_1 \in V_{i_1}$ and $n - t_2 \in V_{i_2}$. Let $d = gcd(t_1, t_2)$. By Proposition 3, $n - t_1 = k_1(t_1 + t_2) + dj_1 - i_1$ and $n - t_2 = k_2(t_1 + t_2) + dj_2 - i_2$ for some k_1, k_2, j_1 and j_2 such that $0 \leq k_1, k_2 \leq \lfloor \frac{n-1}{t_1+t_2} \rfloor$, $1 \leq j_1, j_2 \leq \frac{t_1+t_2}{d}$ and $0 \leq i_1, i_2 \leq d$. Since $d = gcd(t_1, t_2)$, $i_1 \equiv i_2 \pmod{d}$. Since $0 \leq i_1, i_2 \leq d - 1$, we have $i_1 = i_2$ and $n - t_1$ and $n - t_2$ are in the same component of $T_{n-1}\langle t_1, t_2 \rangle$ and consequently in the same component of G. In the same way, $n - t_1$ and $n - t_i$ are in the same component of G.

Theorem 2 (i) For even k, the Toeplitz graph $G = T_n \langle t_1, \ldots, t_k \rangle$ is r-regular if and only if r = k, $n = t_i + t_{k-i+1}$, for each $i, 1 \le i \le \frac{k}{2}$. (ii) For odd k, the Toeplitz graph $G = T_n \langle t_1, \ldots, t_k \rangle$ is r-regular if and only if r = k, $n = t_i + t_{k-i+1}$, for each $i, 1 \le i \le \lfloor \frac{k}{2} \rfloor$. Then n is even and $t_{\frac{k+1}{2}} = \frac{n}{2}$.

Proof. First, suppose that $T_n(t_1, \ldots, t_k)$ is an *r*-regular graph. Obviously, $r \leq k$. Suppose that r < k. By the definition of Toeplitz graph, $t_{r+1} < n$. So $t_{r+1}+1 \in V(G)$ and also, $\{1, t_{r+1}+1-t_r, \dots, t_{r+1}+1-t_1\} \subseteq N(t_{r+1}+1)$. Therefore, $d(t_{r+1}+1) \ge r+1$, a contradiction. Thus r = k. We claim that $n = t_1 + t_k$. First, assume that $n < t_1 + t_k$. Thus $d(t_1) < r$, a contradiction. Next, suppose that $n > t_1 + t_k$. Since $t_1 + t_k + 1 \le n$, $d(t_1 + 1) \ge r + 1$. Which is a contradiction. Hence $n = t_1 + t_k$. By induction on k, we prove the rest. The theorem is true for k = 2. Now, we show that the assertion holds for k = 3. Clearly, $T_n \langle t_1, t_2, t_3 \rangle = T_n \langle t_1, t_3 \rangle \cup T_n \langle t_2 \rangle$. Note that $T_n\langle t_1, t_3\rangle$ and $T_n\langle t_2\rangle$ are edge disjoint. By Proposition 2, $T_n\langle t_1, t_3\rangle$ is a disjoint union of cycles. Also, $T_n(t_1, t_2, t_3)$ is a 3-regular graph. As a result, $T_n \langle t_2 \rangle$ is a perfect matching. So n is even and $t_2 = \frac{n}{2}$. Suppose that the assertion holds for $T_n(t_1, \ldots, t_s)$, and each s, s < k. It is easy to see that $T_n\langle t_1,\ldots,t_k\rangle = T_n\langle t_1,t_k\rangle \cup T_n\langle t_2,\ldots,t_{k-1}\rangle$ which are edges disjoint. By Proposition 2, $T_n(t_1, t_k)$ is a 2-regular graph. So $T_n(t_2, \ldots, t_{k-1})$ is a (k-2)-regular graph. Now, by induction hypothesis the proof of one side

is complete. For the other side, first suppose that $p \in V(T_n\langle t_1, \ldots, t_k \rangle)$ and $p \leq t_1$. Since $t_1 + t_k = n$, $N(p) = \{p + t_1, \ldots, p + t_k\}$. So d(p) = k. Next, assume that $t_{i-1} , for some <math>i, 1 \leq i \leq k$. Now, we have $N(p) = \{p - t_1, \ldots, p - t_{i-1}, p + t_1, \ldots, p + t_{k-i+1}\}$ because $t_i + t_{k-i+1} = n$. Therefore, for $p, t_{i-1} . Finally, suppose that <math>t_k .$ $Clearly, <math>N(p) = \{p - t_1, \ldots, p - t_k\}$. Thus $T_n\langle t_1, \ldots, t_k\rangle$ is a k-regular graph. \Box

3. The Edge-cut and Vertex-cut in Toeplitz Graphs

In this section, we prove a necessary condition for a Toeplitz graph to be 2-edge connected.

Lemma 4 $T_n \langle t_1, t_2 \rangle$ has no cut vertex if and only if $n \ge t_1 + t_2$.

Proof. Suppose that $n \geq t_1 + t_2$. The proof is by induction on n. If $n = t_1 + t_2$, then by Proposition 2, $T_n \langle t_1, t_2 \rangle$ is the disjoint union of d cycles, where $d = gcd(t_1, t_2)$. Since cycles have no cut vertex, the assertion is true for $n = t_1 + t_2$. Clearly, the graph $T_n \langle t_1, t_2 \rangle$ is constructed by adding the vertex n to $T_{n-1} \langle t_1, t_2 \rangle$ and jointing n to two vertices $n - t_1$ and $n - t_2$. Since $n - 1 \geq t_1 + t_2$, the induction hypothesis shows that $T_{n-1} \langle t_1, t_2 \rangle$ has no cut vertex. By Corollary 3, $n - t_1$ and $n - t_2$ are in the same component of $T_{n-1} \langle t_1, t_2 \rangle$, thus $T_n \langle t_1, t_2 \rangle$ has no cut vertex.

Conversely, suppose that $T_n\langle t_1, t_2 \rangle$ has no cut vertex. On contrary, suppose that $n < t_1 + t_2$ but then $T_n\langle t_1, t_2 \rangle$ will have a cut vertex because by Corollary 2, at least one of its components is a path which completes the proof.

By Lemma 4 and Proposition 3, we have the following corollary.

Corollary 4 If $gcd(t_1, t_2) = 1$ and $t_1 + t_2 \le n$, then $T_n < t_1, t_2 > is$ 2-connected.

Proposition 4 Let $gcd(t_1, t_k) = 1$ and $t_1 + t_k < n$, then $T_n < t_1, \ldots, t_k >$ is 2-connected.

Proof. Since $t_1 + t_k < n$, by Lemma 4, $T_n < t_1, t_k >$ has no cut vertex. Since $gcd(t_1, t_k) = 1$ so $gcd(t_1, \ldots, t_k) = 1$ and $t_1 + t_k < n$ which implies $t_1 + t_k \le n + 1$, by corollary 4, $T_n < t_1, \ldots, t_k >$ is connected. $T_n < t_1, \ldots, t_k >$ has more edges than $T_n < t_1, t_k >$, so $T_n < t_1, \ldots, t_k >$ has no cut vertex, i.e., $T_n < t_1, \ldots, t_k >$ is 2-connected. \Box

Theorem 3 Let n_0 be a positive. If $T_{n_0}\langle t_1, \ldots, t_k \rangle$ has no cut edge, then $T_n\langle t_1, \ldots, t_k \rangle$, $n \geq n_0$, has no cut edge as well.

Proof. The proof is by induction on n. For $n = n_0$, there is nothing to prove. Thus assume that $n > n_0$ and $G = T_n\langle t_1, \ldots, t_k \rangle$. Suppose that the assertion holds for n - 1, and $n > n_0$. By induction hypothesis, $T_{n-1}\langle t_1, \ldots, t_k \rangle$ as a subgraph of $T_n\langle t_1, \ldots, t_k \rangle$ has no cut edge. Therefore, the common edges of $T_{n-1}\langle t_1, \ldots, t_k \rangle$ and $T_n\langle t_1, \ldots, t_k \rangle$ are not cut edge. So it is sufficient to prove that none of the edges $\{n, n - t_i\}, 1 \le i \le k$, is a cut edge. By Corollary 3, $n - t_1, n - t_2, \ldots, n - t_k$ are in the same component of $G = T_n\langle t_1, \ldots, t_k \rangle$ and we are done. \Box

Theorem 4 If $3t_1 + 2t_k \leq n$, then $G = T_n < t_1, \ldots, t_k > has$ no cut edge.

Proof. We show that $H = T_{n-t_1-t_k} < t_1, \ldots, t_k >$ has no cut edge. Assume that $1 \le i \le n - t_1 - t_k$. Clearly,

$$N_H(i) = \{i - t_r | 1 \le r \le k, t_r < i\} \cup \{i + t_r | 1 \le r \le k, i + t_r \le n - t_1 + t_k\}$$

Let $r, 1 \leq r \leq k$. If $t_1 < i - t_r$, then the cycle $(i, i - t_r, i - t_r - t_1, i - t_1, i)$ is a subgraph of $T_{n-t_1-t_k} < t_1, \ldots, t_k >$ which contains the edge $\{i - t_r, i\}$. If $0 < i - t_r \leq t_1$, then $i + t_1 \leq 2t_1 + t_k \leq n - t_1 - t_k$, because $3t_1 + 2t_k \leq n$. Now, the cycle $(i, i - t_r, i - t_r + t_1, i + t_1, i)$ is a subgraph of $T_{n-t_1-t_k} < t_1, \ldots, t_k >$ which contains the edge $\{i - t_r, i\}$. If $i + t_r \leq n - t_1 - t_k$ and $t_1 < i$, then the cycle $(i, i + t_r, i + t_r - t_1, i - t_1, i)$ is a subgraph of $T_{n-t_1-t_k} < t_1, \ldots, t_k >$ which contains the edge $\{i, i + t_r\}$. Since $3t_1 + 2t_k \leq n$, $i + t_r \leq 2t_1 + t_k \leq n - t_1 - t_k$ and $t_1 \geq i$, then the cycle $(i, i + t_r, i + t_r + t_1, i + t_1, i)$ is a subgraph of $T_{n-t_1-t_k} < t_1, \ldots, t_k >$ which contains the edge $\{i, i + t_r\}$. Since $3t_1 + 2t_k \leq n$, $i + t_r \leq 2t_1 + t_k \leq n - t_1 - t_k$. So if $i + t_r \leq n - t_1 - t_k$ and $t_1 \geq i$, then the cycle $(i, i + t_r, i + t_r + t_1, i + t_1, i)$ is a subgraph of $T_{n-t_1-t_k} < t_1, \ldots, t_k >$ which contains the edge $\{i, i + t_r\}$. Thus $T_{n-t_1-t_k} < t_1, \ldots, t_k >$ has no cut edge. Now, by Theorem 3, $G = T_n < t_1, \ldots, t_k >$ has no cut edge and the proof is complete.

The following results were proved about the connectivity of Toeplitz graphs.

Remark 1 There is a Toeplitz graph $T_n < t_1, \ldots, t_k >$ such that $gcd(t_1, \ldots, t_k) =$ 1 and $t_1 + t_k \le n + 1$ which is connected but it is not 2-edge connected. For example $T_7 < 3, 5 >$.

Theorem 5 If $gcd(t_1, t_k) = 1$ and $t_1 + t_k \leq n$, then $T_n < t_1, \ldots, t_k > is$ 2-edge connected.

Proof. By Corollary 4, $T_n < t_1, \ldots, t_k >$ is a connected graph. Since $gcd(t_1, t_k) = 1$, by Proposition 2, $T_{t_1+t_k} < t_1, t_k >$ is a cycle. So $T_{t_1+t_k} < t_1, t_k >$ is 2-edge connected. Now, Theorem 3 shows that $T_n < t_1, t_k >$ has no cut edge. Thus the Toeplitz graph $T_n < t_1, \ldots, t_k >$ is 2-edge connected. \Box

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